



# Type Classes in Haskell

Tom Schrijvers

**KU LEUVEN**

Leuven Haskell  
User Group

# Haskell Research Team



Monads

Type  
Classes

GHC

Folds

Pattern  
Matching

Equational  
Reasoning

DSLs

Advanced  
Types

A large, stylized, semi-transparent 'X' graphic in a dark blue-grey color, centered on the slide. The 'X' is composed of two overlapping 'V' shapes, one pointing up and one pointing down.

Adhoc Overloading

# Example: Addition

Different Implementations for different types

`Int -> Int -> Int`

`Float -> Float -> Float`

`Vector -> Vector -> Vector`

...

How do we name the implementations?

# Different Names for Different Implementations

e.g. Ocaml

`+` : `int -> int -> int`



`+.`  : `float -> float -> float`



- remember more names
- common concept / properties

# Same Name for Different Implementations

e.g. Java

`+ : int -> int -> int`



`+ : float -> float -> float`

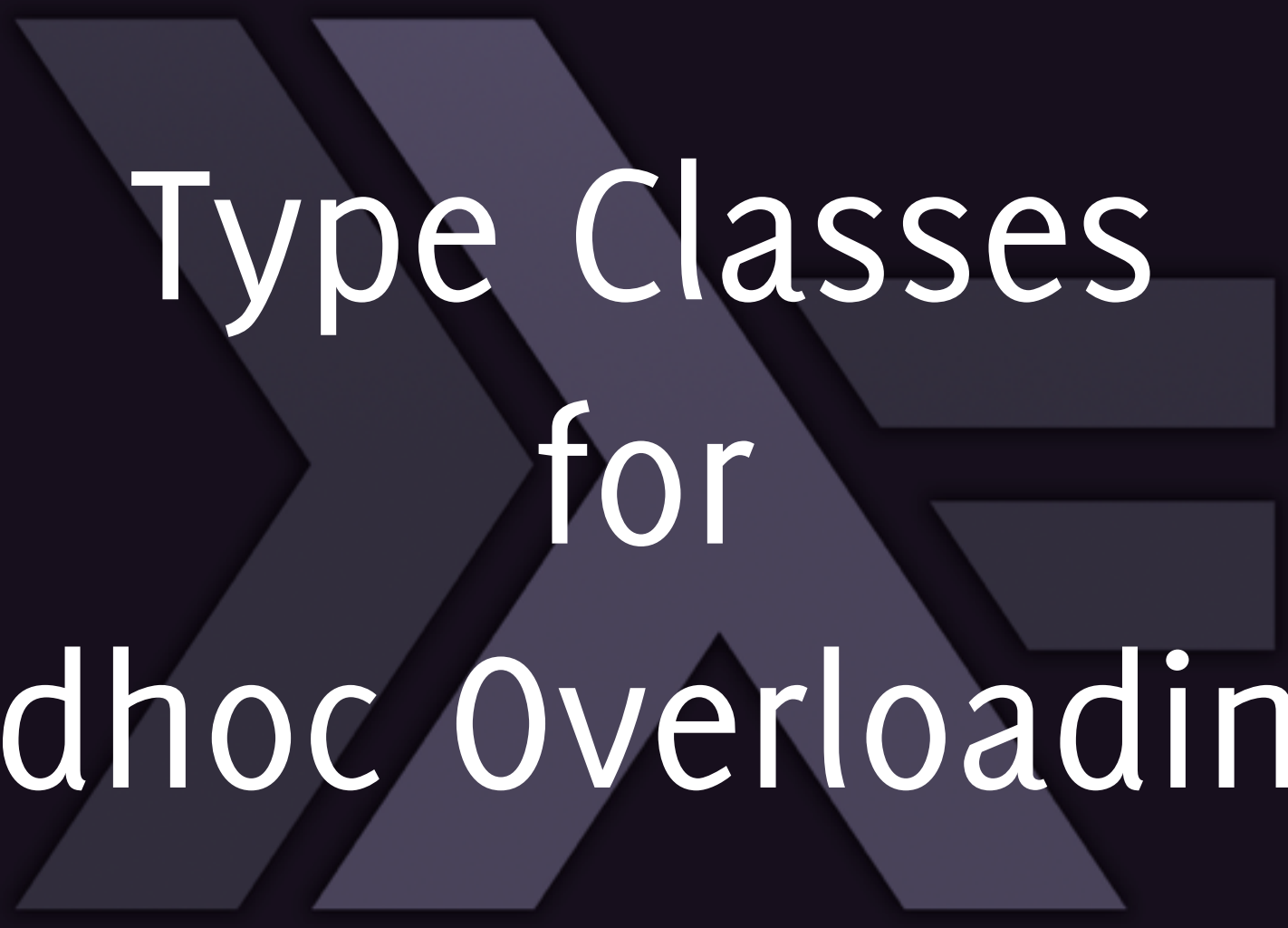
ad hoc overloading

operator overloading

ad hoc polymorphism



- hard-wired in the language
- not user-extensible



# Type Classes for Adhoc Overloading

# Type Classes

“How to make  
ad-hoc polymorphism  
less ad hoc”

Philip Wadler and Stephen Blott. 1989.



# Example: Equality

`(==) :: a -> a -> Bool`

equality is adhoc  
overloaded

```
> "hello" == "world"  
False
```

```
> 2 == (1+1)  
True
```

# Example: Equality

`(==) :: a -> a -> Bool`

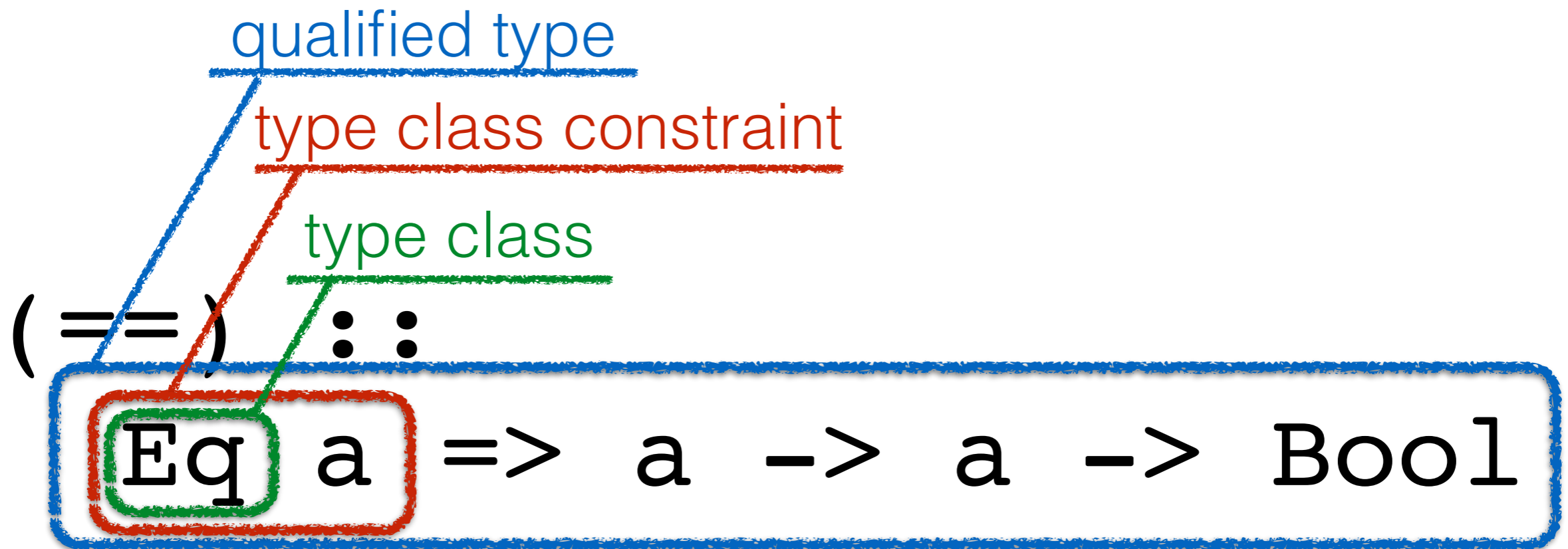
not every type supports equality

```
> id == (\x -> x)
```

No instance for `(Eq (a -> a))`  
arising from a use of `'=='`

In the expression: `id == (\ x -> x)`

# Constraint Polymorphism



type `a` must have an implementation of equality

# Type Class Declaration

```
class Eq a where  
(==) :: a -> a -> Bool
```

method

# Type Class Instantiation

```
data Color = Red | Blue
```

```
> Red == Red
```

```
True
```

```
> Red == Blue
```

```
False
```

# Using Methods

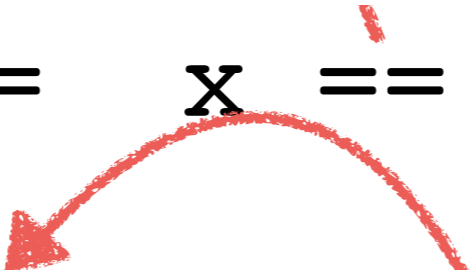
```
isRed :: Color -> Bool  
isRed c = c == Red
```

ok because  
Eq Color instance  
exists

# Constraint Polymorphic Use

most general type?

same  $(x, y) = x == y$



type class constraints propagate!

# Multiple Methods

```
class Eq a where  
  (==) :: a -> a -> Bool  
  (/=) :: a -> a -> Bool
```



# Type Class Instantiation

```
data Color = Red | Blue
```

```
instance Eq Color where
```

```
  Red == Red = True
```

```
  Blue == Blue = True
```

```
  _ == _ = False
```

```
  x /= y = not (x == y)
```

# Default Implementations

```
class Eq a where  
  (==) :: a -> a -> Bool  
  (/=) :: a -> a -> Bool  
  x /= y = not (x == y)
```

# Default Implementation

```
data Color = Red | Blue
```

```
instance Eq Color where
```

```
  Red == Red = True
```

```
  Blue == Blue = True
```

```
  _ == _ = False
```

```
x /= y = not (x == y)
```

# Default Implementations

```
class Eq a where
```

```
  (==) :: a -> a -> Bool
```

```
  x == y = not (x /= y)
```

```
  (/=) :: a -> a -> Bool
```

```
  x /= y = not (x == y)
```

```
class Eq a where
```

The **Eq** class defines equality (**==**) and inequality (**/=**). All the basic datatypes exported by the **Prelude**, and **Eq** may be derived for any datatype whose constituents are also instances of **Eq**.

Minimal complete definition: either **==** or **/=**.

**Minimal complete definition**

(==) | (/=)


# Instance Contexts

instance context

```
instance Eq a => Eq [a] where  
  [] == [] = True  
  (x:xs) == (y:ys) = x==y && xs==ys  
  _ == _ = False
```

# Instance Contexts

```
instance (Eq a, Eq b) => Eq (a, b) where  
  (x1, x2) == (y1, y2)  
    = x1 == y1 && x2 == y2
```



# Subclassing

super class

```
class Eq a => Ord a where  
  (<=) :: a -> a -> Bool
```

every instance of Ord  
must also have  
an instance of Eq

```
r' :: Ord a => a -> a -> Bool  
r' x y = x == y
```

# The Num Class

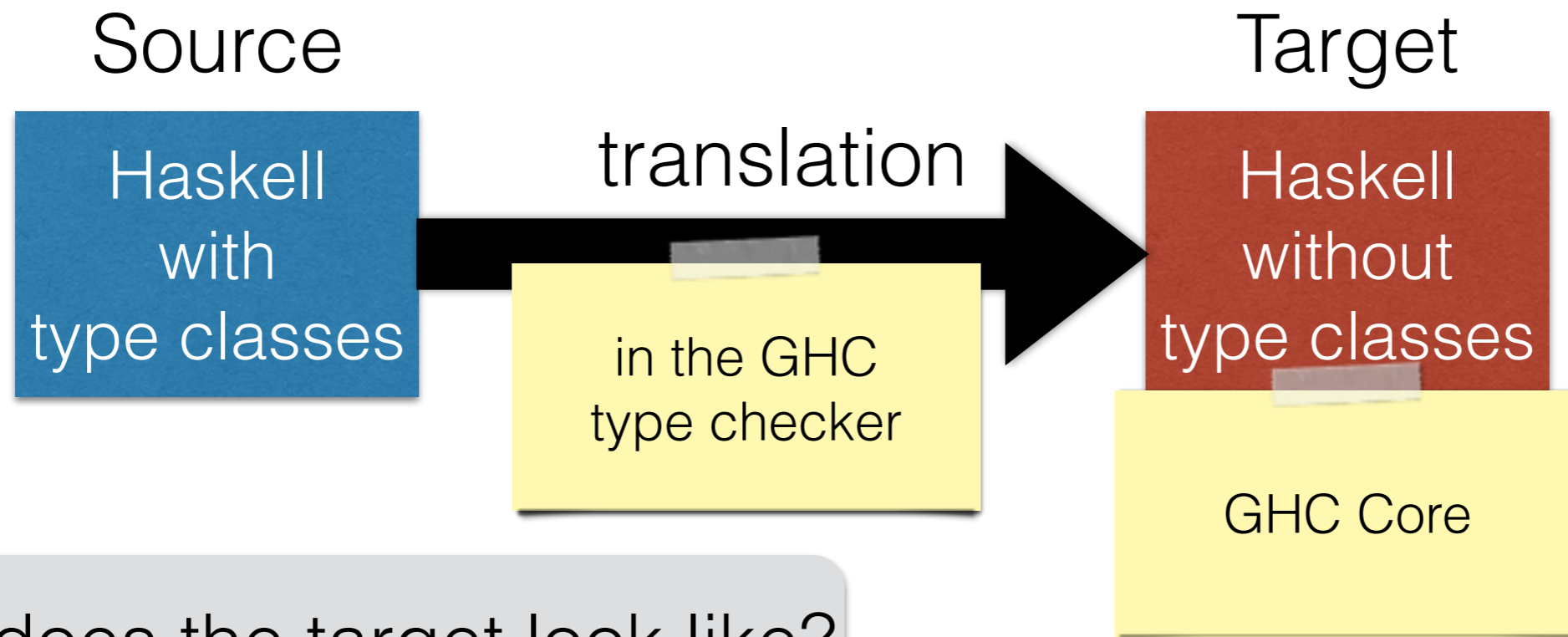
```
class Num a where  
  (+)      :: a -> a -> a  
  (-)      :: a -> a -> a  
  (*)      :: a -> a -> a  
  negate   :: a -> a  
  abs      :: a -> a  
  signum   :: a -> a  
  fromInteger :: Integer -> a
```



A large, stylized, semi-transparent 'X' graphic is centered in the background. It is composed of two overlapping 'X' shapes, one slightly offset from the other, creating a layered effect. The color is a dark, muted purple or blue.

# Dictionary-Passing Translation

# Implementation



What does the target look like?

How is the target obtained?

# Dictionary Representation

```
class Eq a where  
  (==) :: a -> a -> Bool  
  (/=) :: a -> a -> Bool
```



```
data Eq a =  
  Eq { (==) :: a -> a -> Bool  
      , (/=) :: a -> a -> Bool  
      }
```

(function)  
dictionary

# Instance Translation

```
instance Eq () where  
  (==) = \x y -> True  
  (/=) = \x y -> False
```



```
eqUnit :: Eq ()  
eqUnit =  
  Eq { (==) = \x y -> True  
      , (/=) = \x y -> False  
      }
```

# Call Translation

```
p :: () -> () -> Bool  
p x y = x == y
```




```
p :: () -> () -> Bool  
p x y = (==) eqUnit x y
```

typically followed  
by inlining

# Instance Translation


```
instance (Eq a, Eq b) => Eq (a,b) where  
  (==) = \ (x1,y1) (x2,y2)  
        -> x1 == x2 && y1 == y2
```

```
eqPair :: Eq a -> Eq b -> Eq (a,b)  
eqPair da db =  
  Eq { (==) =  
        \ (x1,y1) (x2,y2)  
        -> (==) d1 x1 x2 &&  
            (==) d2 y1 y2  
        , ... }
```



# Call Translation

```
q :: (() , ()) -> (() , ()) -> Bool
q x y = x == y
```




```
p :: (() , ()) -> (() , ()) -> Bool
p x y = (==) d x y
      where
          d = eqPair eqUnit eqUnit
```

# Constraint Polymorphic Signature Translation

```
r :: Eq a => a -> a -> Bool  
r x y = x == y
```


```
r :: Eq a -> a -> a -> Bool  
r d x y = (==) d x y
```





# Constraint Polymorphic Signature Translation


```
s :: Eq a => (a, a) -> (a, a) -> Bool  
s x y = x == y
```



```
s :: Eq a -> (a, a) -> (a, a) -> Bool  
s d x y = (==) d' x y  
  where  
    d' = eqPair d d
```

# Super Class Translation

```
class Eq a => Ord a where  
  (==) :: a -> a -> Bool  
  (/=) :: a -> a -> Bool
```



```
data Ord a =  
  Ord { super :: Eq a  
       / ...  
       }
```

# Super Class Access

```
r' :: Ord a => a -> a -> Bool  
r' x y = x == y
```



```
r' :: Ord a -> a -> a -> Bool  
r' d x y = (==) d' x y  
      where  
          d' = super d
```

Resolution

# Resolution Process

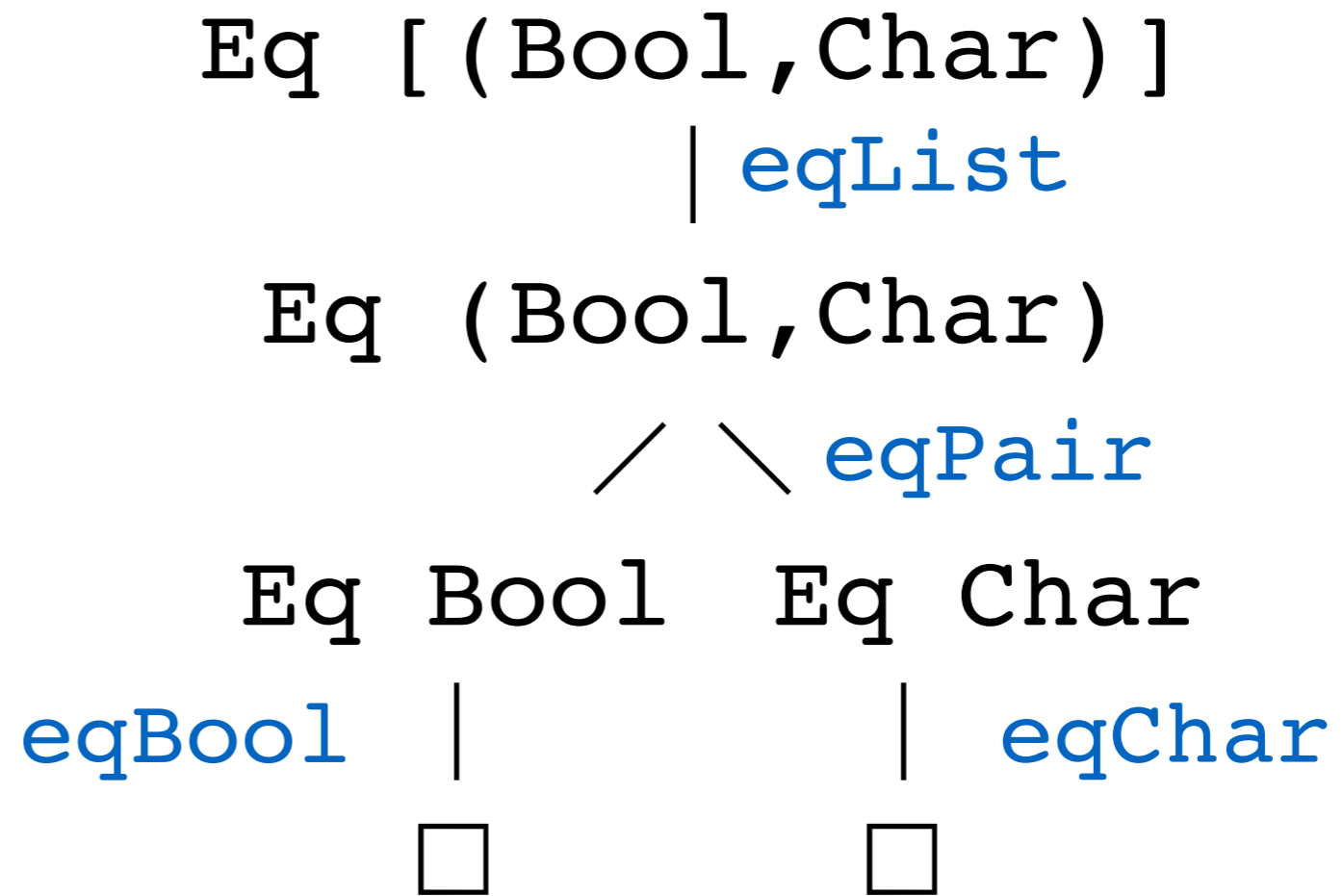
```
[ (True, 'a') ] == [ (False, 'b') ]
```

```
(==) ? [] [ (True, 'a') ] [ (False, 'b') ]
```



```
Eq [ (Bool, Char) ]
```

# Resolution Process



```
instance Eq a ==> Eq [a]
instance (Eq a, Eq b) ==> Eq (a, b)
instance Eq Bool
instance Eq Char
...
```

# Resolution Process

```
(==) d [] [(True, 'a')] [(False, 'b')]
```

```
Eq [(Bool, Char)]
```

```
  | eqList
```

```
Eq (Bool, Char)
```

```
  / \ eqPair
```

```
Eq Bool  Eq Char
```

```
eqBool  |  |  eqChar
```

```
□      □
```

**where**

```
d = eqList (eqPair eqBool eqChar)
```



# Summary



# Type Classes

- ★ **Systematic approach** to adhoc overloading
- ★ implemented by means of **dictionary passing**
- ★ where **resolution** assembles the dictionaries.

# LANGUAGE Pragmas

`MultiParamTypeClasses`

`FunctionalDependencies`

`FlexibleContexts`

`FlexibleInstances`

`UndecidableInstances`

`OverlappingInstances`

`IncoherentInstances`

non-terminating  
resolution possible

incoherent  
resolution possible

incoherent  
resolution likely

Join the Google Group

Leuven Haskell User Group

Join the Meetup

Leuven Haskell User Group